

# EECS2011 Fundamentals of Data Structures (Winter 2022)

Q&A - Week 2 Lecture

Wednesday, January 26

## Announcements

- Lecture W3 released (SLL) 
- - Assignment 1 (requiring SLL) to be released on Monday.
- Plan of Returning In-Person (starting Feb. 14)
  - + Unchanged
    - \* Pre-recorded lectures
    - \* Zoom Weekly Q&A and Office hours in the first instance
    - \* Online Programming & Written tests in the first instance
  - + Changed
    - \* In-Person Exam
  - + To be determined:
    - \* Some (programming and/or written) tests may be in-person, in which case you'll be notified at least one week in advance.

# Inserting into an Array



```
String[] insertAt(String[] a, int n, String e, int i)
    String[] result = new String[n + 1];
    for(int j = 0; j <= i - 1; j ++){ result[j] = a[j]; }
    result[i] = e;
    for(int j = i + 1; j <= n; j ++){ result[j] = a[j-1]; }
    return result;
```

For the second loop, what's the worst case of the number of iterations?

$$\begin{array}{ccccccc} \cdot & \cdot & \cdot & & & & \\ - & [i+1, n] & & n - (i+1) + 1 & = & \boxed{n-i} & \rightarrow \# \text{ iterations} \\ - & \text{insertAt}(a, a.length, e, 0) & & & & & \end{array}$$

to maximize this #

↓  
 $\bar{i}$  is minimized

↓  
 $\bar{i} = 0$ .

$$[z, n-1] \rightarrow (n-1) - z + 1 = \underline{\underline{n-z}}$$

for (int  $i=0$ ;  $i < n$ ;  $i++$ ) {

~~for (int  $j=i$ ;  $j < n$ ;  $j++$ ) {~~

~~$\dots$~~   ~~$j$ :  $n$~~

$O(1)$

$\dots$

$\underline{i}$   
 $\underline{0} =$   
 $\underline{1} =$   
 $\underline{2} =$   
- - -  
 $\underline{n-1}$

$\underline{j}$   
 $\underline{0} \quad 1 \quad 2 \dots (n-1)$   
→  $(n-1)$  Iterations  
 $\underline{1} \quad 2 \dots (n-1)$   
→  $(n-2)$  Iterations  
 $\underline{2} \dots (n-1)$

$i$  It.  
 $\underline{(n-1)}$

$$\begin{aligned} & \underline{n} + (n-1) + (n-2) + \dots + \underline{1} \\ &= \frac{(n+1) \cdot n}{2} \end{aligned}$$

$O(n^2)$

Thursday,  
Jan. 20  
O&R

for ( $i = n ; i \geq 1 ;$ ) {

$O(1)$



$i = \underline{i/2} ;$   
3

$$n = \underline{1024}$$

$\underline{i}$

1024

512

256

128

64

32

16

8

4

2

1

iteration

$i$	Iteration
1024	1
512	2
256	3
128	4
64	5
32	6
16	7
8	8
4	9
2	10
1	11

$\log_2 1024$   
"10".

## Alternative Solution to groupSum

Credit: Mohammad P.

```
public boolean groupSum(int start, int[] nums, int target) {  
    if (start >= nums.length) {  
        return target == 0;  
    }  
    else {  
        return groupSum(start+1, nums, target-nums[start]) ||  
            groupSum(start+1, nums, target);  
    }  
}
```

# Problem on Recursion: Climbing Staircase

Climb  
steps

You are asked to program this method:

Staircase

**public int climb(int h, int n)**

The return value suggests the number of possible ways  
for climbing a stair of height h, while each climb is up to n steps.

Assumptions:  $n \leq h$ , each climb takes at least 1 step

e.g.,  $\text{climb}(4, 2)$  returns 5, meaning that there are 5 ways  
for climbing a stair of height 4 and each climb takes up to 2 steps.

1. 1 step, 1 step, 1 step
2. 1 step, 1 step, 2 steps.
3. 1 step, 2 steps, 1 step.
4. 2 steps, 1 step, 1 step
5. 2 steps, 2 steps.

first climb:  
1 step      first climb: 2 steps

$\text{climb}(4-1, 2)$        $1, 2, \dots, n$   
"      3       $\text{climb}(h-n, n)$

$\text{climb}(4-2, 2) \rightarrow$  1 step,  $\text{climb}(h-1, n)$   
"      2 steps,  $\text{climb}(h-2, n)$   
"       $i$  steps,  $\text{climb}(h-i, n)$

strictly  
smaller prob.

